

Stochastic Dynamic Teams and Games with Asymmetric Information

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OUTLINE

- Introduction to decision making in stochastic environments
- Dynamic teams, games, role of information, asymmetry
- Some caveats and counter-examples on existence and characterization of sols
- Existence and characterization of team-optimal policies under asymmetry
- Existence and characterization of Nash equilibrium policies under asymmetry
- Conclusions

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General Framework

INGREDIENTS

- *Uncertainty* (uncertain environment)
- *Decision makers* (DMs) (players, agents)
- *Perceptions* of DMs on uncertainty
- *Objective(s)*
- Description of interactions (underlying *network*)
- Common information / Private information

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General Framework

- *Uncertainty* (uncertain environment); *decision makers* (DMs) (players, agents); *perceptions* of DMs on uncertainty; *objective(s)*; description of interactions (underlying *network*); common information
- DMs pick *policies* (decision laws, strategies) leading to *actions* that evolve over time

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- Policies are constructed based on information received (active as well as passive) and guided by individual utility or cost functions over the DM horizon
- Single DM => stochastic control
- Single objective => stochastic teams
- Otherwise ZS or NZS games, with NE

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Coupling of Information and Actions

- Is the quality of active and relevant information received by a player affected by actions of other players?
 - If no => underlying game (team) is generally "simple"
 - If yes => it is generally "difficult"

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Asymmetric Information

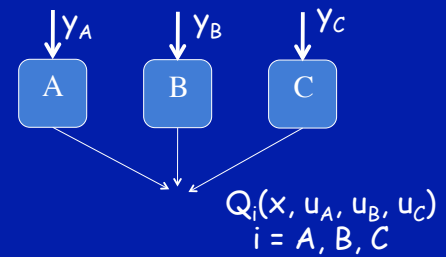
A stochastic decision problem is one with *asymmetric information*, if different decision units (players, agents, decision makers) acting at the same time instant do not have access to the same information (of relevance).

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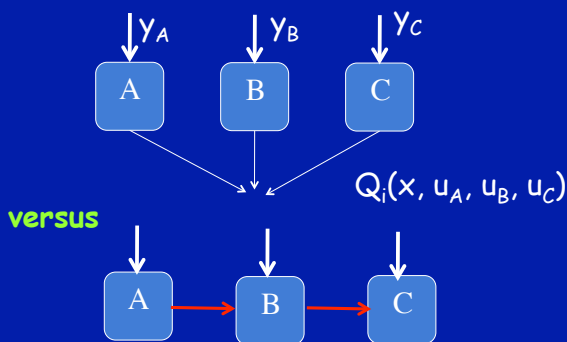
Non-classical Information

A stochastic decision problem is one with *non-classical information*, if a decision unit, **B**, that "follows" another one, **A**, does not have all the information acquired and used by **A**.

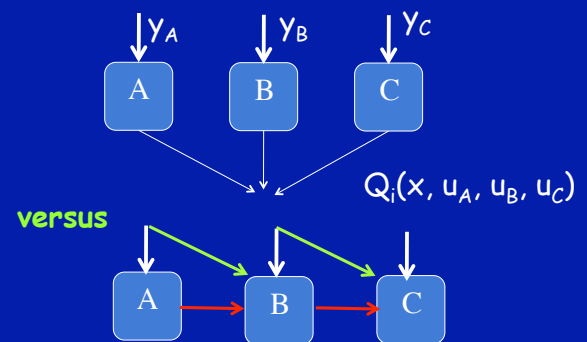
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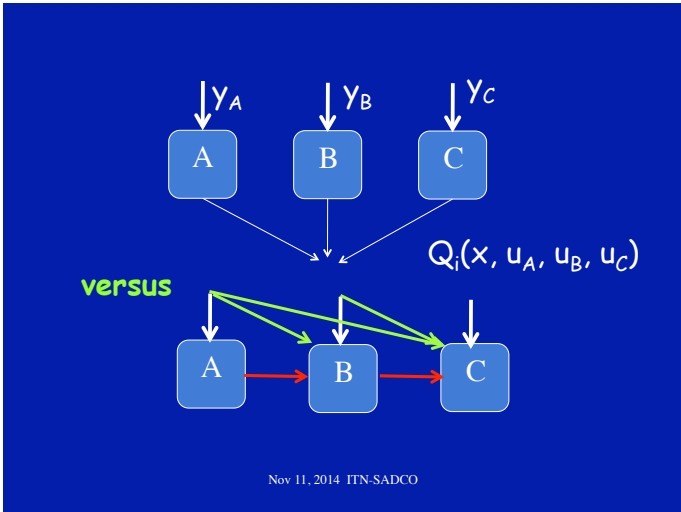
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- ### Questions / Challenges
- Existence of team-optimal (T-O) solutions
 - Existence of Nash equilibrium (NE) policies (in dynamic stochastic games)
 - Characterization of T-O solutions
 - Characterization of NE policies
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 - Characterization of T-O solutions
 - Characterization of NE policies
 - **Challenges**
 - Possibility of non-existence with asymmetry
 - Triple role (caution, probing, signaling)
 - Signaling (deception/threat) through action
 - Tension between signaling and optimization
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- ### Questions / Challenges
- NEXT: An example exhibiting the subtleties, such as the significance of failure of continuity of information content of channels in the limit of a sequence of discretized problems*
- **Challenges**
 - Possibility of non-existence with asymmetry
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SCP with non-classical information (limited memory)

$x, w \sim$ independent random variables
 u_0, u_1 are real valued actions
 $J(\gamma_0, \gamma_1) = E [Q(x, u_0, u_1) \mid \gamma_0, \gamma_1]$
 $J^* = \min \min J(\gamma_0, \gamma_1)$

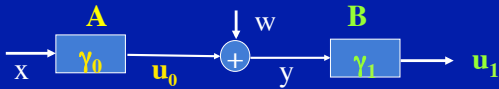
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Non-classical SCP (2-pt pmf's)

$x \sim \pm 1$ with equal prob; likewise w
 $J(\gamma_0, \gamma_1) = E [k (u_0 - x)^2 + (u_0 - u_1)^2 \mid \gamma_0, \gamma_1], k > 0$
 $J^* = \inf \inf J(\gamma_0, \gamma_1)$

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Infirmizing policies



$J^* = \inf \inf J(\gamma_0, \gamma_1) = 0$ - lowest possible
Consider the policies:

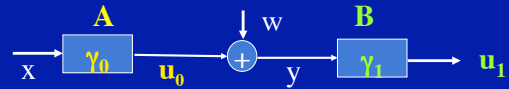
$$\gamma_0(x) = x + \varepsilon \operatorname{sgn}(x)$$

$$\gamma_1(y) = 1 + \varepsilon, \text{ if } y = 2 + \varepsilon \text{ or } \varepsilon \\ = -1 - \varepsilon, \text{ if } y = -2 - \varepsilon \text{ or } -\varepsilon$$

$$\rightarrow J = k \varepsilon^2 \rightarrow 0 \text{ as } \varepsilon \rightarrow 0$$

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Infirmizing policies



But, as $\varepsilon \rightarrow 0$, limiting policies do not lead to $J = 0$!

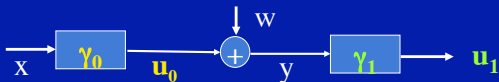
Loss of crucial information in the limit!

$$\gamma_1(y) = 1 + \varepsilon, \text{ if } y = 2 + \varepsilon \text{ or } \varepsilon \\ = -1 - \varepsilon, \text{ if } y = -2 - \varepsilon \text{ or } -\varepsilon$$

$$\rightarrow J = k \varepsilon^2 \rightarrow 0 \text{ as } \varepsilon \rightarrow 0$$

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Same SCP but with Gaussian RVs (Witsenhausen, 1968)



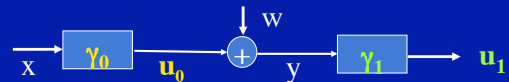
$$x \sim N(0, \sigma_x^2) \quad w \sim N(0, \sigma_w^2)$$

$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

→ optimal team solution exists, but its structure is not known

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Same SCP but with Gaussian RVs (Witsenhausen, 1968)

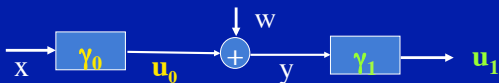


$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

→ optimal team solution exists, but its structure is not known
affine policies are not optimal

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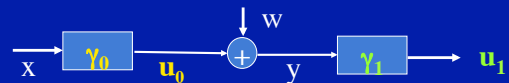


$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

Note that if $u_0 = \gamma_0(x)$, best γ_1 is
 $u_1 = \gamma_1(y) = E[\gamma_0(x) | y]$
Hence optimization is with respect to the probability distribution of u_0

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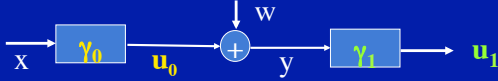


$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

Original existence proof given by Witsenhausen involved calculus of variations type arguments, and does not extend to more general scenarios.

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A class of non-linear policies



$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

A policy pair that beats the best linear one:

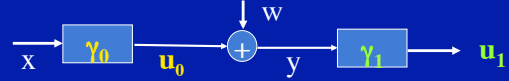
$$u_0 = \gamma_0(x) = \varepsilon \operatorname{sgn}(x) + \lambda x$$

$$u_1 = \gamma_1(y) = E[\varepsilon \operatorname{sgn}(x) + \lambda x \mid y]$$

optimize wrt ε and λ

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A class of non-linear policies



$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

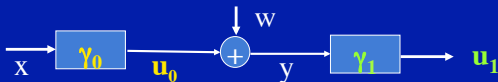
A more general one with improved perf:

$$u_0 = \gamma_0(x) = \varepsilon \operatorname{Quant}(x) + \lambda x$$

$$u_1 = \gamma_1(y) = E[\varepsilon \operatorname{Quant}(x) + \lambda x \mid y]$$

optimize wrt ε and λ for different **Quant(ization)** schemes (Bansal, TB '87)

A class of non-linear policies



$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

$$u_0 = \gamma_0(x) = \varepsilon \operatorname{Quant}(x) + \lambda x$$

$$u_1 = \gamma_1(y) = E[\varepsilon \operatorname{Quant}(x) + \lambda x \mid y]$$

Best perf under linear could be arbitrarily *bad* against best performance under **Quant**

$$\Rightarrow \sup_{k, \sigma} [Q_W^{\text{linear}^*} / Q_W^{\text{Quant}^*}] = \infty \quad (\text{Mitter, Sahai '99})$$

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BUT

Is the solution to Wits (68) piecewise affine ?

NO

It is strictly increasing with a real analytic left inverse

(Wu, Verdú - CDC'11)

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NO

It is strictly increasing with a real analytic left inverse

(Wu, Verdú CDC'11)

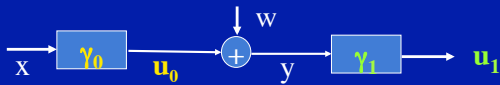
Uses optimal transport theory:

Given prob measures P and Q , and cost function $c: \mathbb{R}^2 \rightarrow \mathbb{R}$, find inf over all joint distributions of X and Y , with marginals P and Q , of $E[c(X, Y)]$

(Monge-Kantorovich)

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Alternative existence proof

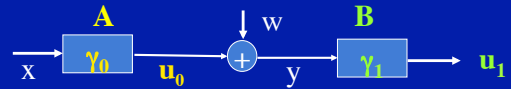


$$Q_W(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

$E[c(x, u_0)] = E[Q_W] = k W_2(P, Q)^2 + \text{mmse}(Q)$
 where P and Q are PDF's of x and u_0 , and W_2 is the quadratic Wasserstein distance between P and Q .

For any P , \exists minimizing Q , with "smooth" γ_0 .

Earlier general set-up with a different Q



$$x \sim N(0, \sigma_x^2) \quad w \sim N(0, \sigma_w^2)$$

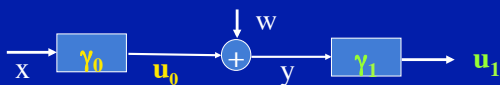
u_0, u_1 are real valued actions

$$J(\gamma_0, \gamma_1) = E[Q(x, u_0, u_1) | \gamma_0, \gamma_1]$$

$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2$$

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Gaussian Test Channel

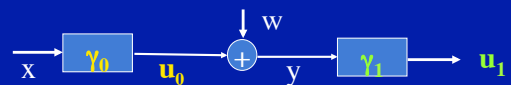


$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2$$

→ optimal control law (encoder/decoder) exists, and is linear

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Generalized Gaussian Test Channel

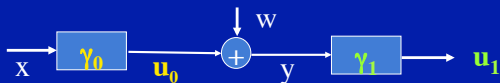


$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

→ optimal control law (encoder/decoder) exists, and is linear

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Generalized Gaussian Test Channel (Proof of existence)



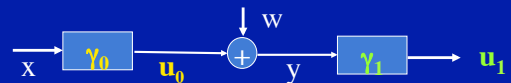
$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

$$E[Q] = F(\gamma_0, \gamma_1) \geq k\alpha + \beta + \inf_{\gamma} b_0 E[\gamma_0(x)x]$$

$$\geq k\alpha + \beta - \text{sgn}(b_0) \sigma_x \sqrt{\alpha}$$

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Generalized Gaussian Test Channel (Proof of existence-2)



$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

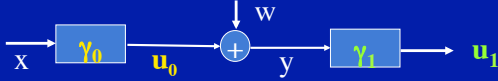
$$E[Q] = F(\gamma_0, \gamma_1) \geq k\alpha + \beta + \inf_{\gamma} b_0 E[\gamma_0(x)x]$$

$$\geq k\alpha + \beta - \text{sgn}(b_0) \sigma_x \sqrt{\alpha}$$

DPT: $I(X; Y) \geq I(X; U_1)$

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Generalized Gaussian Test Channel (Proof of existence-3)



$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

$$E[Q] = F(\gamma_0, \gamma_1) \geq k\alpha + \beta + \inf_{\gamma} b_0 E[\gamma_0(x)x]$$

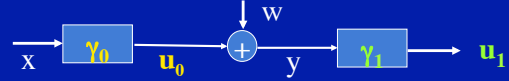
$$\geq k\alpha + \beta - \text{sgn}(b_0) \sigma_x \sqrt{\alpha}$$

$$(1/2) \log(1 + (\alpha/\sigma_w^2)) \geq I(X;Y) \geq I(X;U_1) \geq (1/2) \log(\sigma_x^2/\beta)$$

$C(\alpha)$ $R(\beta)$

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Generalized Gaussian Test Channel (Proof of existence-4)



$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

$$E[Q] = F(\gamma_0, \gamma_1) \geq k\alpha + \beta + \inf_{\gamma} b_0 E[\gamma_0(x)x]$$

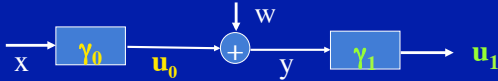
$$\geq k\alpha + \beta - \text{sgn}(b_0) \sigma_x \sqrt{\alpha}$$

$$(1/2) \log(1 + (\alpha/\sigma_w^2)) \geq I(X;Y) \geq I(X;U_1) \geq (1/2) \log(\sigma_x^2/\beta)$$

$$\Rightarrow \beta \geq \sigma_w^2 \sigma_x^2 / (\sigma_w^2 + \alpha)$$

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Generalized Gaussian Test Channel (Proof of existence-5)



$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

$$E[Q] = F(\gamma_0, \gamma_1) \geq k\alpha + \beta + \inf_{\gamma} b_0 E[\gamma_0(x)x]$$

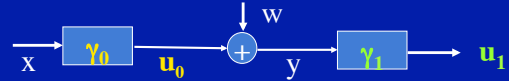
$$\geq k\alpha + \beta - \text{sgn}(b_0) \sigma_x \sqrt{\alpha}$$

$$\Rightarrow \beta \geq \sigma_w^2 \sigma_x^2 / (\sigma_w^2 + \alpha)$$

Inequality is tight with $\gamma_0(x) = -\text{sgn}(b_0)(\sqrt{\alpha}/\sigma_x) x$

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Generalized Gaussian Test Channel (Proof of existence-6)



$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

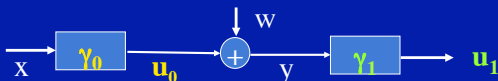
$$E[Q] = F(\gamma_0, \gamma_1) \geq k\alpha + \beta - |b_0| \sigma_x \sqrt{\alpha}$$

$$\geq k\alpha + \sigma_w^2 \sigma_x^2 / (\sigma_w^2 + \alpha) - |b_0| \sigma_x \sqrt{\alpha}$$

Obtain the α that minimizes the bound $\rightarrow \alpha^*$
Then, $\gamma_0^*(x) = -\text{sgn}(b_0)(\sqrt{\alpha^*}/\sigma_x) x$, $\gamma_1^*(y) = E[x|y]$

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Generalized Gaussian Test Channel



$$Q(x, u_0, u_1) = k(u_0)^2 + (u_1 - x)^2 + b_0 u_0 x$$

One of the few instances when static/causal coding (and linear in this case) leads to attainment of equality in $C(\alpha) \geq R(\beta)$

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Relay Channels Multiple Serial Decision Units



$$x \sim N(0, \sigma_x^2), w_i \sim N(0, \sigma_w^2), v \sim N(0, \sigma_v^2)$$

$$J(\gamma_0, \gamma_{[1,m]}) = E[Q(x, u_0, u_{[1,m]}) | \gamma_0, \gamma_{[1,m]}]$$

$$Q(x, u_0, u_{[1,m]}) = (u_m - x)^2 + \sum_{i=1}^m (u_{i-1})^2$$

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Multiple Serial Decision Units



$$x \sim N(0, \sigma_x^2), \quad w_i \sim N(0, \sigma_w^2), \quad v \sim N(0, \sigma_v^2)$$

$$J(\gamma_0, \gamma_{[1,m]}) = E [Q(x, u_0, u_{[1,m]}) \mid \gamma_0, \gamma_{[1,m]}]$$

→ For $m > 1$, affine laws no longer optimal for this extended GTC model

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Does a general unifying existence proof exist?

One that would apply to

- Witsenhausen counter-example
- Gaussian test channel
- their multi-dimensional versions
- relay channels
- stochastic control problems with no memory or limited memory
- LQG teams with asymmetric information
- etc

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Answer is YES

(Gupta, Yüksel, TB -- CDC'14)

Assumptions on the dynamic M -agent T stage stochastic team problem:

- Sequential decision problem: Agent i observes and recalls at time t , $\gamma_{(i,t)}$, and constructs $u_{(i,t)} = \gamma_{(i,t)}(\gamma_{(i,t)})$
- $\inf J(y)$ is finite (e.g. use zero control)
- Witsenhausen's static reduction holds

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Answer is YES

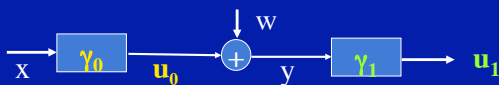
(Gupta, Yüksel, TB -- CDC'14)

Static reduction:

- Convert M -agent team to MT -agent one where each agent acts only once
- Measure and cost transformation that turns the dynamic problem into static one, with independent measurements (measurements now enter into cost)

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Static Reduction for the C-example



$$x \sim N(0, \sigma_x^2) \quad w \sim N(0, \sigma_w^2)$$

$$Q_w(x, u_0, u_1) = k(u_0 - x)^2 + (u_0 - u_1)^2$$

→ $\int P(dx)P(dy) Q_w(x, u_0, u_1) \cdot \exp(u_0(2y - u_0)/2)$

$$x \sim N(0, \sigma_x^2) \quad y \sim N(0, 1)$$

independent

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Challenges in the Proof

- Extracting a convergent subsequence of the infimizing sequence
- Showing lower semi-continuity of J
- Making sure that informational constraints are preserved in the limit

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Approach to Overcome Challenges in the Proof

- With static reduction, redefine optimization over randomized strategies
- Because of independence of measurements information constraints are preserved
- Identify a compact set that contains the optimal solution (Markov's inequality)
- Use Blackwell's *irrelevant information theorem* to go to deterministic policies
- Invoke equivalence of reduced static and dynamic teams

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Assumptions & Approach Apply to Various Settings

- Multi-dimensional Witsenhausen C-E
- Multi-dimensional Gaussian Test Channel
- Relay channel and multi-dimensional version
- LQG with static output feedback (solution is generally nonlinear)
- Countable/quantized observation spaces in static teams with observation sharing IS (uses a *lifting* technique)

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Details in: A. Gupta, S. Yüksel & TB, "Existence of optimal strategies in a class of dynamic stochastic teams," CDC 2014

Also: Gupta, Yüksel, Langbort & TB (ACC'14)

And: Gupta, Yüksel, TB & Langbort, "On the existence of optimal policies in a class of sequential dynamic teams," 2014 (arxiv)

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Comprehensive coverage of analysis, optimization and performance of stochastic networked systems in

Stochastic Networked Control Systems: Stabilization and Optimization under Information Constraints -- Yüksel & TB
Birkhäuser, 2013

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Stochastic Dynamic Games with Asymmetric Information

"Lifting" to a symmetric game utilizing the common information of the players

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Formulation

- **Game G1**: n player game with state & measurement for P_i :

$$x_{t+1}^i = f_t^i(x_t^i, x_t, u_t^i, w_t^i)$$

$$y_t^i = h_t^i(x_t^i, x_t, v_t^i), \quad i = 1, \dots, n$$
- Some sharing of past measurements and control actions by the players $\rightarrow I_{(i,t)}$
- Player i 's policy at time t : $g_{(i,t)}: I_{(i,t)} \rightarrow U_{(i,t)}$
- $J^i(g^1, \dots, g^n) = E[\sum_{t=1}^T c_t^i(x_t^i, x_t, u_t^i) | g]$
- Interested in Nash equilibrium g^*

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Transformation to a game with symmetric information (say, $n=2$)

- Decompose into common and private information: $C_t = I_{(1,t)} \wedge I_{(2,t)}$, $P_{(i,t)} = I_{(i,t)} \setminus C_t$
 - Replace P_i with F_i who has access to C_t , and selects $\Gamma_t^i: P_{(i,t)} \rightarrow U^i$ according to φ_t^i (his strategy): $\Gamma_t^i = \varphi_t^i(C_t)$
 - Then, control action is: $u_t^i = \Gamma_t^i(P_{(i,t)})$
 - Cost: $L^i(\varphi^1, \dots, \varphi^n) = E[\sum_{[1, T]} c^i(x_t^i, x_t, \underline{u}_t) | \underline{\varphi}]$
- **Game G2**

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Game G2: Replace P_i with F_i who has access to C_t , and selects $\Gamma_t^i: P_{(i,t)} \rightarrow U^i$ according to φ_t^i (his strategy): $\Gamma_t^i = \varphi_t^i(C_t)$
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Theorem: Let \underline{g} be a NE for **G1**. Define $\varphi_t^i(C_t) = g_t^i(\cdot, C_t)$
Then, $\underline{\varphi}$ is a NE for **G2**. Conversely, if φ is a NE for **G2**, \underline{g} as above is a NE for **G1**.

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What does this lead to (conceptually)

- Since **G2** is a symmetric full information game, its NE can be obtained by backward induction (strongly-time consistent, sub-game perfect, etc), albeit with optimization on a function space $\{\Gamma_t^i\}$
- Markov perfect equilibrium (MPE) - generally one of many other NE which however are not STC

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- NEED some conditions for this to go through

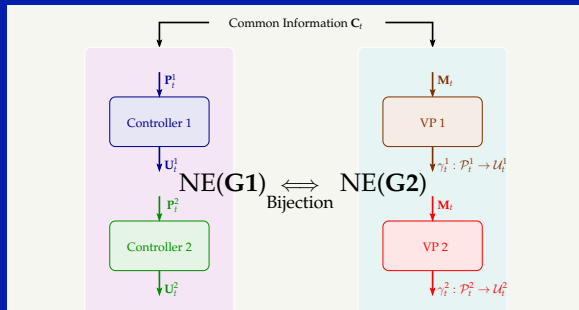
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Assumptions / Approach

- Common information C_t does not contract
- Common information based common belief is independent of the policies of the players
- -----
- Introduced common information based game (**G2**) on lifted strategy spaces
- MPE of **G2** is CIMPE for **G1**

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Relationship between $G1$ & $G2$



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Assumptions / Approach / Result

- Common information C_+ does not contract
- Common information based common belief is independent of the policies of the players
- -----
- Introduced common information based game ($G2$) on lifted strategy spaces as before
- MPE of $G2$ is CIMPE for $G1$
- \rightarrow For LQG, CIMPE is unique & CIMPE strategies are affine; computation involves solving linear eqs

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Recent papers:

Nayyar & TB (CDC, 2012)
"Dynamic Stochastic Games with Asymmetric Information"

Nayyar, Gupta, Langbort, TB (TAC, 59(3), 2014)
"Common information based Markov perfect equilibria for stochastic games with asymmetric information: Finite games"

Gupta, Nayyar, Langbort, TB (SICON, 52(5), 2014)
"Common information based Markov perfect equilibria for Linear-Gaussian games with asymmetric information"

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Recap

- We now have a general theory of existence of team-optimal solutions in stochastic dynamic teams with asymmetric and non-classical IS
- Caveats in computations based on discretization/quantization
- Lifting of SGs with asymmetric information to ones with symmetric information at the expense of increase in complexity
- Common information based MPE offers computational advantages

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THANKS !

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